# Boundary effects in some stochastic geometric models

Stochastic Geometry Days 2024

Xiaochuan Yang (Brunel)

() joint work with

Frankie Higgs and Mathew Penrose (University of Bath)

## Largest nearest neighbour link

Given a point cloud  ${\mathcal X}$  in a metric space, the LNNL tells us how far the most isolated point is from the others

- $ullet \ L_n = \max_{x \in \mathcal{X}} dist(x, \mathcal{X} \setminus \{x\})$
- $ig| ullet \ L_n = \inf\{r: \deg(x) \geq 1, \ orall x \in G(\mathcal{X},r)\}$
- $ullet \ L_n = \inf\{r: \mathcal{X}(B(x,r)) \geq 2, \ orall x \in \mathcal{X}\}$

similarly  $L_{n,k}$ 

#### Full coverage threshold

 $R_{n,k} = \inf\{r: \mathcal{X}(B(x,r)) \geq k, \ orall x \in A\}$ 

- $\bullet \quad L_n \leq R_{n,2}$
- $\bullet \,\, L_{n,k} \leq R_{n,k+1}$

## **Connectivity threshold**

(i) A graph is k-connected, denoted  $\mathcal{K}_k$ , if the removal of k-1 vertices does not disconnect the graph.

$$M_{n,k} = \inf\{r: G(\mathcal{X},r) \in \mathcal{K}_k\}$$

X-X X-X

 $L_{n,k} \leq M_{n,k}$ 

## Result

 $\mathcal{Q}$  Let  $\mathcal{X} = \mathcal{P}_n$ . Under conditions on A, we have the strong law  $rac{nL_n^d}{\log n}\sim rac{nM_n^d}{\log n}\sim rac{nR_n^d}{\log n}\sim c$ where c depends on the geometry of A.

- binomial process
- non-uniform

- general k = k(n) Penrose, PTRF'22, book
  - Penrose, Y, Higgs, '23

# Finite convex polytope

Let  $k=k(n)\sim eta\log(n)$  with  $eta\in[0,\infty].$  Then

$$c = \max_{\phi \in \Phi^*(A)} rac{\hat{H}_eta(D(\phi)/d)}{f_\phi 
ho_\phi}$$

- $\hat{H}_0(x)=x, \hat{H}_\infty(x)=1$  (chg scale to k(n) when  $eta=\infty$ )
- $\Phi^*(A)$ : faces of A, including  $A^o$  viewed as d-dim face
- $D(\phi)$  is the dimension of the face  $\phi$
- $f_{\phi} = \inf_{x \in \phi} f(x)$  where f is density
- $ho_{\phi} = |B(o,1) \cap \kappa_{\phi}|$  angular volume of face  $\phi$

# **Proof: granulation**

Want to show some threshold  $\sim r_n$ 

- draw grid of spacing  $arepsilon r_n$
- random point "activates" closest grid point
- count relevant grid patterns

let's implement this twice, 1 for torus, 2 for metric space

# $nL_n^d \geq c\log n$ : torus

- assume  $\{L_n \leq r_n\}$  happens with  $nr_n^d = c\log n$
- at each  $arepsilon_1 r_n$  lattice point, EITHER non-activated OR at least two points within  $(1+2\sqrt{d}\,arepsilon_1)r_n$

 $\mathbb{P}[.\cup..] = 1 - \mathbb{P}[\{..\}^c | \{..\}^c] \mathbb{P}[\{..\}^c] \leq 1 - \eta \mathbb{P}[\{..\}^c]$ 

• keeps lattices distant  $3r_n$  from each other, by independence

$$(1{-}\dots)^{c_1r_n^{-d}}=O(\exp(-n^{1- heta_dc})) o 0$$

provided  $c < 1/ heta_d$ 

$$nL_n^d \geq c\log n$$
: metric space

- a portion of A is r-packed with packing number  $\Omega(r^{-b})$
- $ig egin{array}{ll} ig egin{array}{ll} \mu(B(x,r)) \leq ar^d & ext{for every } x & ext{in that portion} \end{array}$

$$(1-\eta(nar^d)e^{-nar^d})^{c_1r_n^{-b}}=O(\exp(-n^{b/d-ac}))
ightarrow 0$$

provided c < (b/d)/a

- the ultimate lower bound is the max of these (b/d)/a for all kinds of portions

# $nM_n^d \leq c\log n$ : torus

suppose  $\{M_n>r_n\}$  with  $nr_n^d=c\log n$ , then  $\sharp$  clusters  $\geq 2$ , so one of the following happens

- $\exists$  isolated point
- $\exists$  non-isolated small clusters,  $0 < diam \leq Kr_n$
- ullet  $\exists$  two big clusters,  $diam > Kr_n$

 $\mathbb{P}[\exists iso] \leq r_n^{-d} \exp(-c heta_d\log n) = n^{1-c heta_d} o 0$  provided  $c > 1/ heta_d.$ 

# $nM_n^d \leq c\log n$ : metric space

- a portion of A is covered by  $O(r^{-b})$  balls with radius r and  $\mu(B(x,r)) \geq ar^d$  for each x in that portion
- $\bullet \ \mathbb{P}[\exists iso \ \text{in portion}] \leq O(r_n^{-b}) \exp(-nar_n^d) = n^{b/d-ac} \rightarrow 0$

provided c > (b/d)/a.

- union bound:  $c > \max(b/d)/a$  over all portions
- small and big clusters are more complicated, draw

 $\bigcirc$  provided that A has matching covering and packing exponent and one can estimate volume of balls ( $G_r\setminus G$ ) optimally (up to epsilon), all three thresholds are asymptotically the same

**Isolated points and AB coverage** Given  $\mathcal{X}, \mathcal{Y}$  sampled from A

$$T = \inf\{r: \mathcal{X}(B(y,r)) \geq 1, orall y \in \mathcal{Y}\}$$

- lower bound of  $R_n$ , the minimal matching threshold, and connectivity threshold of BRGG
- related to AB-percolation, Iyer-Yogeshwaren

 $n\mathbb{E}[|x\in A:\mathcal{P}_n(B(x,r))=0|]=\int_A \exp(-n heta_d r^d)ndx=\mathbb{E}[x]$ 

 $\left( \mathcal{Y}(\mathrm{uncovered}_{r_n}) 
ightarrow Po(.\,) ext{ for some } r_n \Rightarrow T_n \stackrel{d}{
ightarrow} (2compt)Gumbel 
ight)$ 

## Papers

- Largest nearest-neighbour link and connectivity threshold in a polytopal random sample *Journal of Applied and Computational Topology*
- Covering one point process with another *submitted*

code:

https://github.com/frankiehiggs/connectivity-in-polytopes https://github.com/frankiehiggs/CovXY package: https://github.com/xiaochuany/geography

#### Bonne anniversaire, Anne!